

Effective Concurrency with Algebraic Effects

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Concurrency \neq Parallelism

- Concurrency
 - Programming technique
 - **Overlapped** execution of processes
- Parallelism
 - Performance hack
 - **Simultaneous** execution of computations

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(Fibers) (Domains)

Schedulers

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 - Data parallel —> Gang scheduling
- *Algebraic Effects and Handlers*

Algebraic Effects and Handlers

- Programming and reasoning about computational effects in a pure setting.
 - Cf. Monads
- Effects in practice
 - M Pretnar, A Bauer, “Eff programming language”
 - <http://www.eff-lang.org/>
 - O Kiselyov, A Sabry, C Swords, B Foppa, “Extensible-effects for Haskell”
 - <https://hackage.haskell.org/package/extensible-effects>
 - E Brady, “Effects in Idris”
 - <http://eb.host.cs.st-andrews.ac.uk/drafts/eff-tutorial.pdf>
 - O Kammar, S Lindley, N Oury , “Handlers in Action”, ICFP ’13
 - dl.acm.org/citation.cfm?id=2500590

Algebraic Effects: Example

```
exception Foo of int
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```
let f () = 1 + (raise (Foo 3))
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```
let r =  
  try  
    f ()  
  with Foo i -> i + 1
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```
val r : int = 4
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effect Foo : int -> int
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```
let f () = 1 + (perform (Foo 3))
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```
let r =  
  try  
    f () (* spawned in a new fiber *)  
  with effect (Foo i) k ->  
    continue k (i + 1)
```

```
val r : int = 5
```

Effects interface

```
type _ eff += Foo : int -> int eff
```

```
val perform : 'a eff -> 'a
```

```
type ('a, 'b) continuation
```

```
val continue : ('a, 'b) continuation -> 'a -> 'b
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Handlers are Deep!

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val r : int = 8
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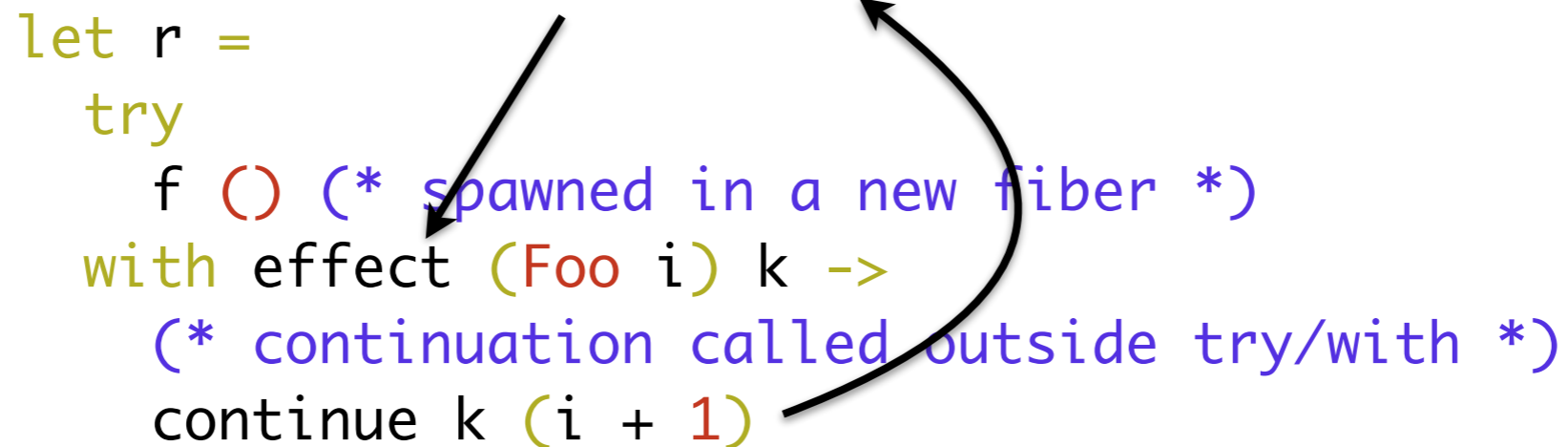
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A diagram consisting of two black arrows. One arrow starts at the function call `f ()` inside the `try` block and points to the `(Foo 3)` argument in the `perform` call within the `let f` definition. The second arrow starts at the `continue k (i + 1)` line in the `with effect` block and points back to the `f ()` call in the `try` block, illustrating the return path from the handler to the caller.

```
val r : int = 8
```

Scheduler Demo¹

[1] <https://github.com/kaycesrk/ocaml15-eff/tree/master/chameneos-redux>

Implementation

- Fibers: Heap allocated, dynamically resized stacks
 - ~10s of bytes
 - No unnecessary closure allocation costs unlike CPS

Implementation

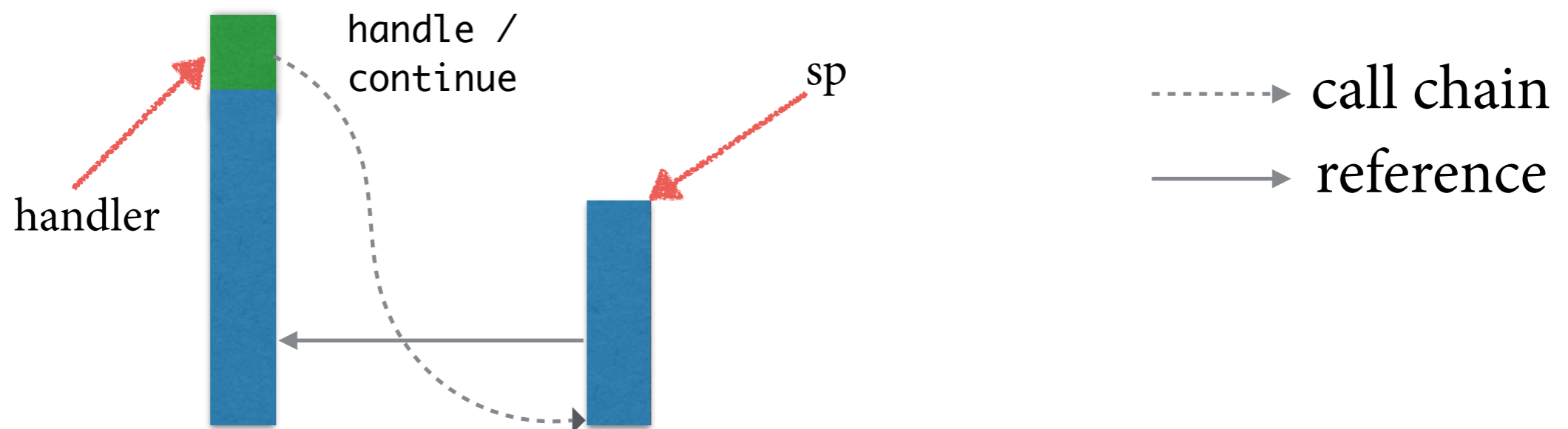
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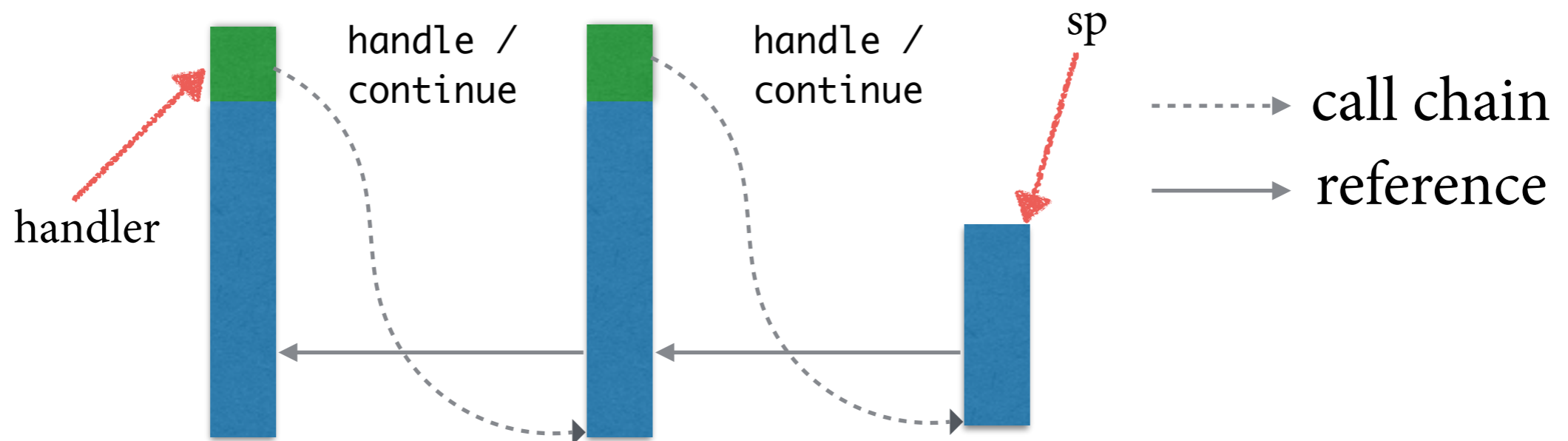
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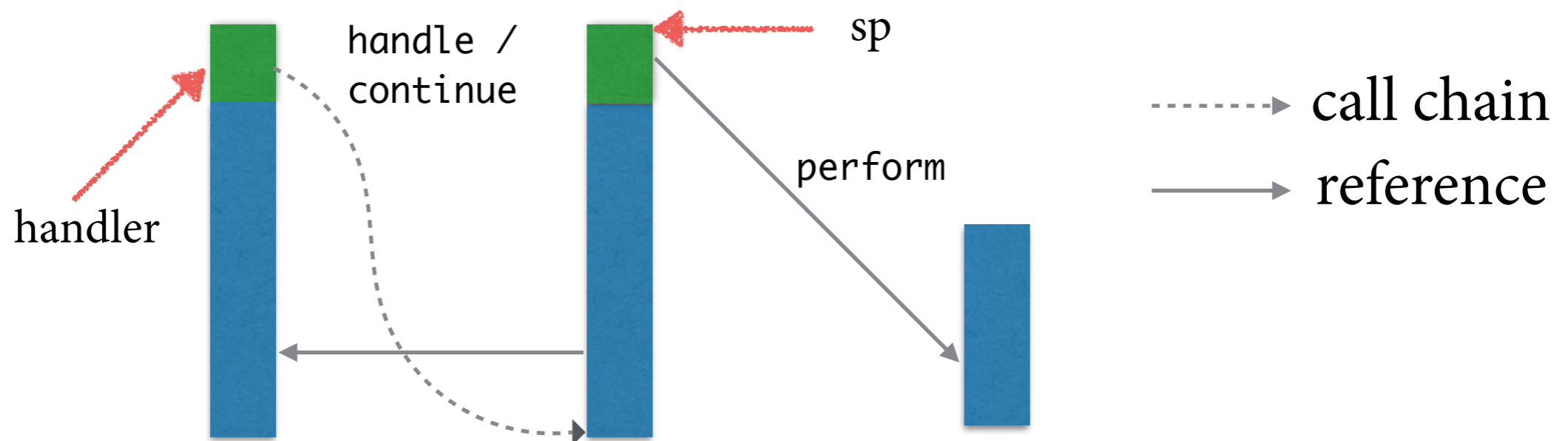
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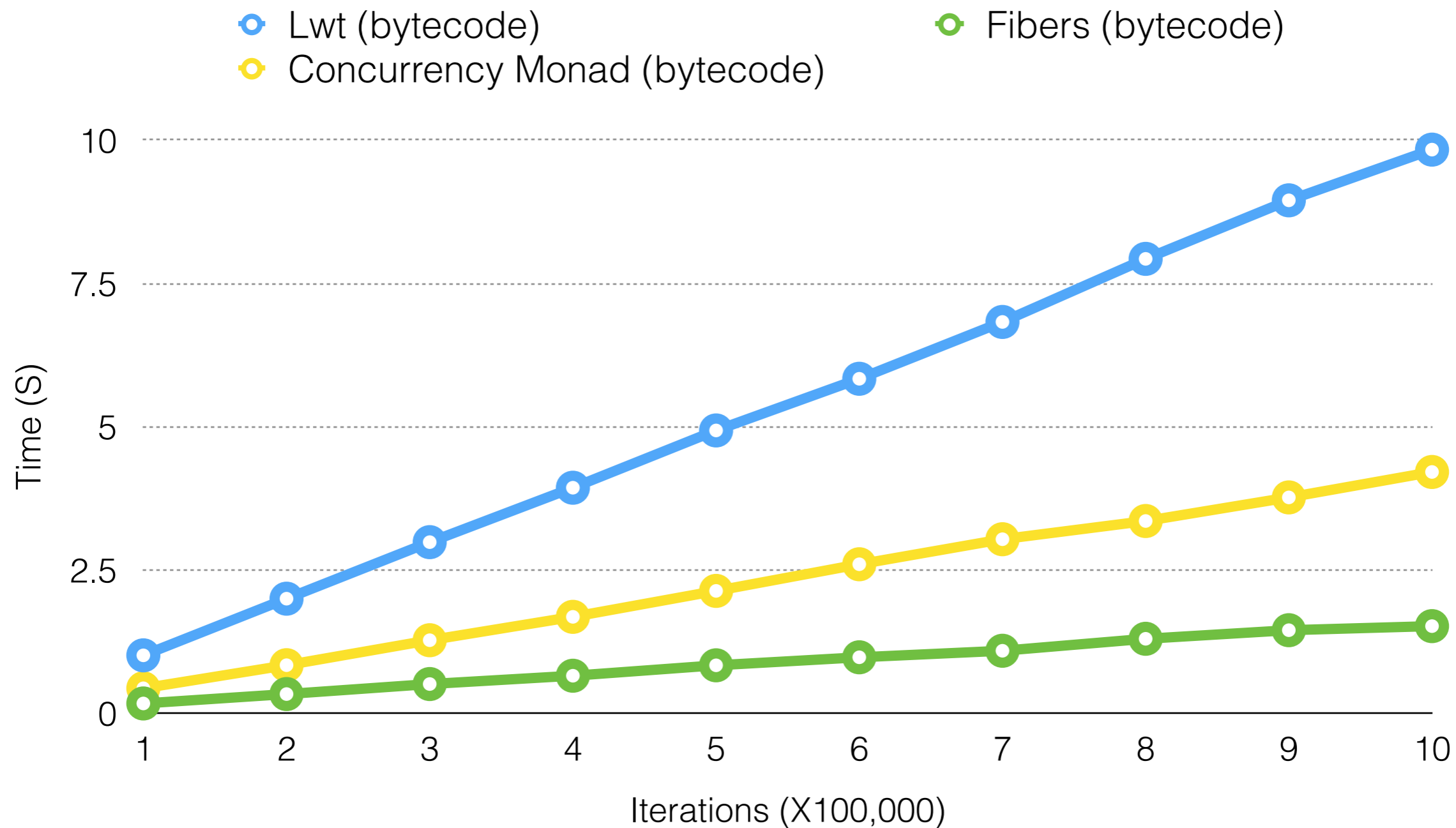


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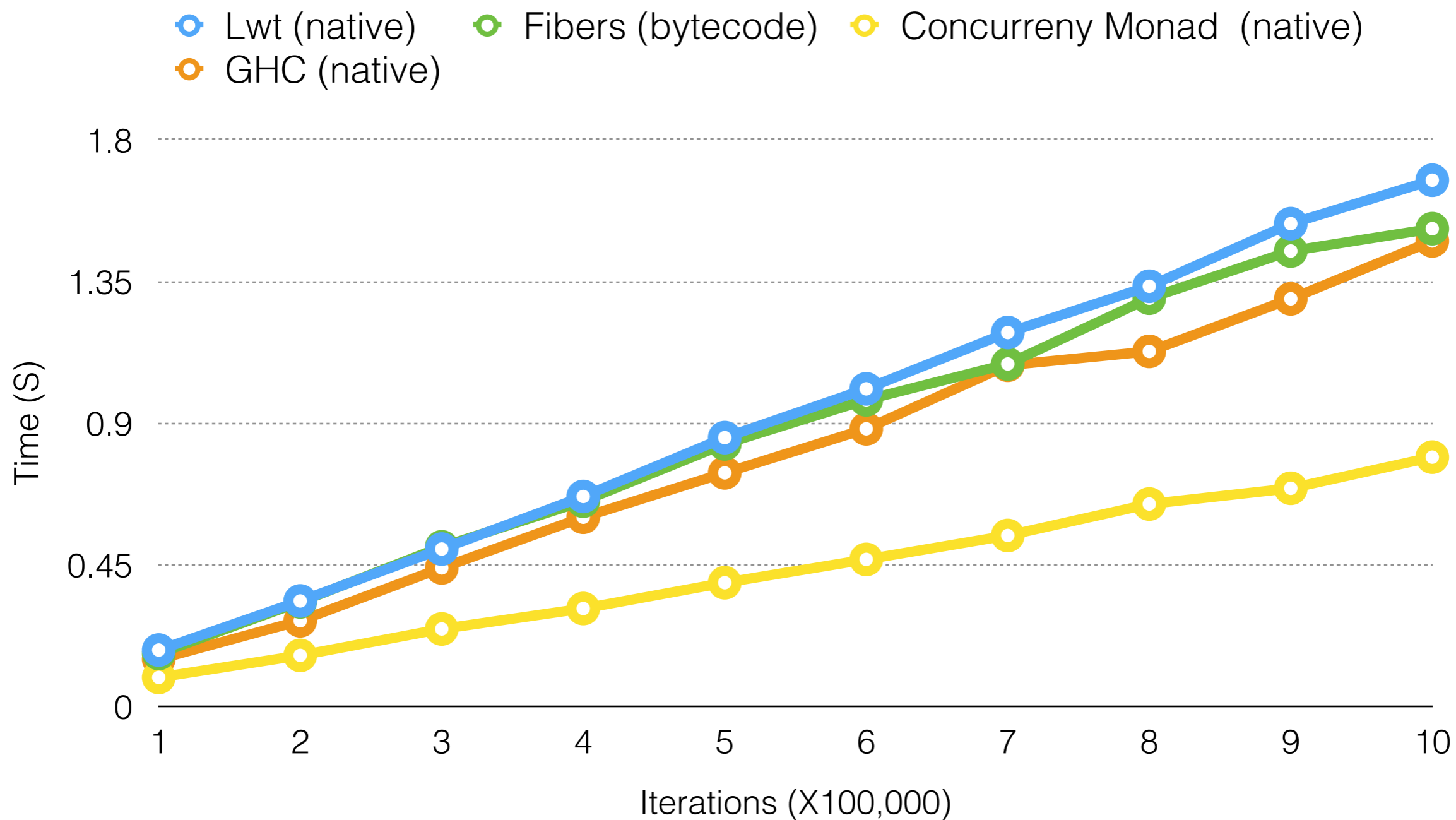
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Performance : Chameneos-Redux



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Generator from Iterator¹

```
type 'a t =  
  | Leaf  
  | Node of 'a t * 'a * 'a t  
  
let rec iter f = function  
  | Leaf -> ()  
  | Node (l, x, r) -> iter f l; f x; iter f r
```

[1] <https://github.com/kayceesrk/ocaml15-eff/blob/master/generator.ml>

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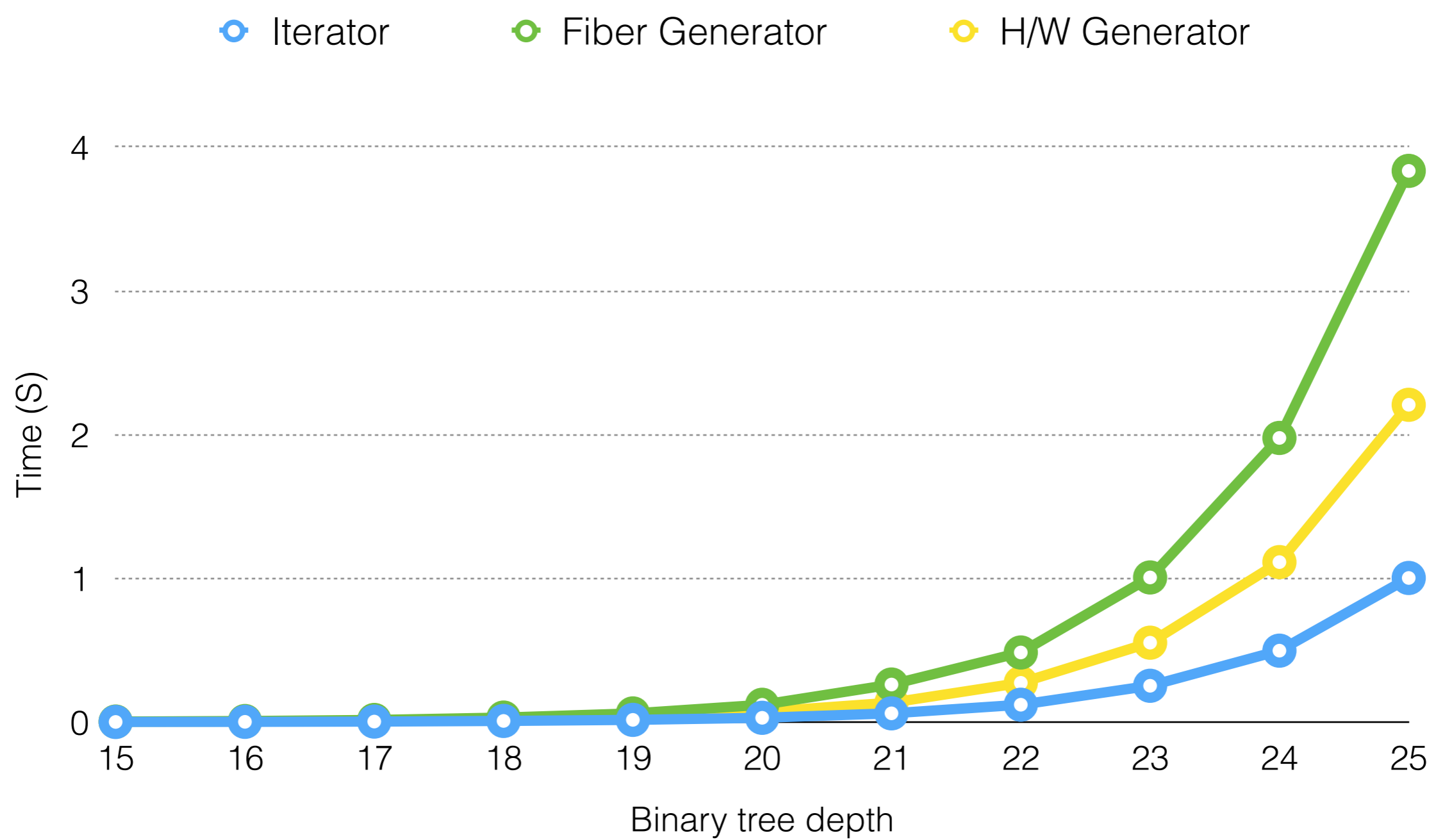
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| Leaf
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let rec iter f = function
| Leaf -> ()
| Node (l, x, r) -> iter f l; f x; iter f r

(* val to_gen : 'a t -> (unit -> 'a option) *)
let to_gen (type a) (t : a t) =
  let module M = struct effect Next : a -> unit end in
  let open M in
  let step = ref (fun () -> assert false) in
  let first_step () =
    try
      iter (fun x -> perform (Next x)) t; None
    with effect (Next v) k ->
      step := continue k; Some v
  in
  step := first_step;
  fun () -> !step ()
```

[1] <https://github.com/kayceesrk/ocaml15-eff/blob/master/generator.ml>

Performance : Generator



Concerns

- Unchecked effects
 - Risks \sim exceptions
 - Effect inference in Eff^1

[1] Matija Pretnar, “Inferring Algebraic Effects”, <http://arxiv.org/abs/1312.2334>

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 - Effect inference in Eff¹
- Interfacing with monadic code (Lwt, Async)
 - Use monadic reflection to recover direct-style code²

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[2] https://github.com/kayceesrk/ocaml15-eff/blob/master/reify_reflect.ml

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- Unchecked effects
 - Risks \sim exceptions
 - Effect inference in Eff¹
- Interfacing with monadic code (Lwt, Async)
 - Use monadic reflection to recover direct-style code²
- Compilation to other backends (JS, Java?)
 - ES6 generators, ES7 async/await
 - Selective-CPS transform³

[1] Matija Pretnar, “Inferring Algebraic Effects”, <http://arxiv.org/abs/1312.2334>

[2] https://github.com/kayceesrk/ocaml15-eff/blob/master/reify_reflect.ml

[3] T Rompf et al., “Implementing first-class polymorphic delimited continuations by a type-directed selective CPS-transform”, ICFP ‘09

Status

- Bytecode only. *Todo Native.*
- Several opportunities for optimisation
 - Continuations invoked at tail position
 - Dynamic search for effect handler
- Code
 - Multicore OCaml: <https://github.com/ocaml-labs/ocaml-multicore>
 - Stand-alone effects: <https://github.com/kayceesrk/ocaml/tree/effects>
 - Effects examples: <https://github.com/kayceesrk/ocaml-eff-example>