

OCaml 5.0

"KC" Sivaramakrishnan







What's in the can?

OCaml 5.0

Merge Process







FAQs



Moving to OCaml 5.0



Concurrency

Parallelism

Concurrency



Overlapped execution

Effect Handlers

Parallelism

Concurrency



Overlapped execution

Effect Handlers

Parallelism



Simultaneous execution

Domains



Domains

- Units of parallelism
- Heavy-weight entities
 - Recommended to have I domain per core



- Non-blocking synchronisation Atomic
- Domain-local state

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- Heavy-weight entities
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- ✦ Create and destroy Spawn and Join
- Blocking synchronisation Mutex, Condition and Semaphore





Domains with Threads







Blocking and non-blocking synchronisation works uniformly across threads and domains



Domainslib

• A library for nested-parallel programming (OpenMP, Cilk, NESL,...)



Pool O





let next () =

```
for x = 0 to board_size - 1 do
  for y = 0 to board_size - 1 do
   next_board.(x).(y) <- next_cell cur_board x y</pre>
  done
done;
```

. . .

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 done
done;
```

. . .

let next () =

T.parallel_for pool ~start:0 ~finish:(board_size - 1) \sim body:(fun x -> for y = 0 to board_size - 1 do next_board.(x).(y) <- next_cell cur_board x y</pre> done);

. . .



Performance: Game of Life

Board size = 1024, Iterations = 512

Cores	Time (Seconds)	Vs Serial
1	24.326	1
2	12.290	1.980
4	6.260	3.890
8	3.238	7.51
16	1.726	14.09
24	1.212	20.07

Allocation and Collection



Domain 0

Domain I

- Minor heap allocations require no synchronization
- Major heap allocator is
 - **Small:** Thread-local, size-segmented free list ◆
 - Large: malloc •

Allocation and Collection



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Domain I

- Minor heap allocations require *no* synchronization
- Major heap allocator is
 - **Small:** Thread-local, size-segmented free list ◆
 - Large: malloc
- Goal is to match *best-fit* for sequential programs
 - + If we're slower than best-fit, then it is a performance regression

Concurrent GC





Start of major cycle

- Stop-the-world parallel *minor* GC + non-moving *major* GC
 - Objects don't move while the mutator is running!



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No additional rules for the C FFI in OCaml 5.0

Same rules as OCaml 4.x hold even for parallel programs! •

- Simple (comprehensible!) operational memory model
 - Only atomic and non-atomic locations
 - DRF-SC
 - No "out of thin air" values
 - Squeeze at most perf \Rightarrow write that module in C, C++ or Rust.



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- Key innovation: Local data race freedom
 - Permits compositional reasoning



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 - + DRF-SC
 - No "out of thin air" values
 - Squeeze at most perf \Rightarrow write that module in C, C++ or Rust.
- Key innovation: Local data race freedom
 - Permits compositional reasoning
- Performance impact
 - Free on x86 and < 1% on ARM</p>



• Simple (comprehensible!) operational memory model

Bounding Data Races in Space and Time

(Extended version, with appendices)

Stephen Dolan University of Cambridge, UK

KC Sivaramakrishnan University of Cambridge, UK

Abstract The primary reasoning tools provided to programmers by these models are the *data-race-freedom* (DRF) theorems. Pro-We propose a new semantics for shared-memory parallel grammers are required to mark as *atomic* all variables used programs that gives strong guarantees even in the presence for synchronisation between threads, and to avoid *data races*, of data races. Our local data race freedom property guarwhich are concurrent accesses (except concurrent reads) to antees that all data-race-free portions of programs exhibit

• Performance impact

+ Free on x86 and < 1% on ARM

PLDI'18

Anil Madhavapeddy University of Cambridge, UK

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- OCaml 5.0 extended the work to cover
 - Object initialisation
 - Compilation to CII memory model
- C FFI has been made stronger (by making the access volatile)

#define Field(x, i) (((volatile value *)(x)) [I])

void caml_modify (volatile value *, value);

void caml_initialize (volatile value *, value);

- Assumes Linux Kernel Memory Model (LKMM)
- Does not break code

• C FFI also respects LDRF!

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let msg = ref 0 let flag = Atomic.make false

let t1 =msg := 1; Atomic.set flag true

let $t_2 =$ let rf = Atomic.get flag in let rm = !msg in assert (not (rf = true && rm = 0))

• C FFI also respects LDRF!

let msg = ref 0 let flag = Atomic.make false

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let $t_2 =$ let rf = Atomic.get flag in let rm = !msg in assert (not (rf = true && rm = 0))

/* t1 implemented in C */

void t1 (value msg, value flag) { caml_modify (&Field(msg,0), Val_int(1)); caml_atomic_exchange (flag, Val_true); }

ThreadSanitizer

```
type t = { mutable x : int }
1
2
   let v = { x = 0 }
3
4
   let () =
5
6
7
     Domain.join t1;
8
     Domain.join t2
9
```

let t1 = Domain.spawn (fun () -> v.x <- 10) in</pre> let t2 = Domain.spawn (fun () -> Unix.sleep v.x) in

ThreadSanitizer

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type t = { mutable x : int }
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   let v = \{ x = 0 \}
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4
   let () =
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     let t1 = Domain.spawn (fun () -> v.x <- 10) in
6
     let t2 = Domain.spawn (fun () -> Unix.sleep v.x) in
7
     Domain.join t1;
8
     Domain.join t2
9
```

WARNING: ThreadSanitizer: data race (pid=502344) Read of size 8 at 0x7fc0b15fe458 by thread T4 (mutexes: write M0): #0 camIDune __exe __Simple_race __fun_600 /workspace_root/simple_race.ml:7 (simple_race.exe+0x51e9b1) #1 caml_callback ??:? (simple_race.exe+0x5777f0) #2 domain_thread_func domain.c:? (simple_race.exe+0x57b8fc)

Previous write of size 8 at 0x7fc0b15fe458 by thread T1 (mutexes: write M1): #0 camIDune __exe __Simple_race __fun_596 /workspace_root/simple_race.ml:6 (simple_race.exe+0x51e971) #1 caml_callback ??:? (simple_race.exe+0x5777f0) #2 domain_thread_func domain.c:? (simple_race.exe+0x57b8fc)

- Structured programming with delimited continuations
- No effect system, no dedicated syntax
- Provides both <u>deep</u> and <u>shallow</u> handlers

Effect handlers

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Example prints "0 I 2 3 4"

Effect handlers

```
open Effect
 1
    open Effect.Deep
2
 3
    type _ Effect.t += E : string t
 4
 5
     let comp() =
 6
      print_string "0 ";
 7
      print_string (perform E);
 8
      print_string "3 "
 9
10
     let main () =
11
      try_with comp ()
12
      { effc = fun (type a) (e : a t) ->
13
14
           match e with
           | E -> Some (fun (k : (a,_) continuation) ->
15
               print_string "1 ";
16
               continue k "2 ";
17
               print_string "4 ")
18
             e -> None }
19
```



- Structured programming with delimited continuations
- No effect system, no dedicated syntax
- Provides both <u>deep</u> and <u>shallow</u> handlers

Example prints "0 I 2 3 4"

• Same type safety as the earlier syntactic version

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             e -> None }
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```



Eio — Direct-style effect-based concurrency



HTTP server performance using 24 cores



HTTP server scaling maintaining a constant load of 1.5 million requests per second

Integration with Lwt & Async



- Lwt_eio allows running Lwt and Eio code together
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 - Incrementally port Lwt applications to Eio

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- Lwt_eio allows running Lwt and Eio code together
 - Only sequential
 - Cancellation semantics is also integrated
 - Incrementally port Lwt applications to Eio
- Very experimental ASYNC_ei0 running Async and Eio code together
 - Required changes to Async

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 - Multiple tricky rebases to keep the fork up to date with trunk
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 - Not worth splitting into multiple PR context loss
- Asynchronous & Synchronous review phases (Nov 2021)

Working Group	Multicore Lead	Dev Team Lead
GC	Tom Kelly & Sadiq Jaffer	Damien Doligez
Domains	Tom Kelly	Luc Maranget
Runtime multi-domain safety	Enguerrand Decorne	Xavier Leroy
Stdlib changes	KC Sivaramakrishnan	Florian Angeletti (and Gabriel Scherer)
Fibers	KC Sivaramakrishnan	Damien Doligez (and Xavier Leroy)

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<> Cod	e 💽 Issues 268 រ៉ ្ Pull requests 258 🕞	Actions 🗄 Projects	() Security	🗠 Insights		
Merged Edit <> Code - Image: Second state of the second state of t						
	kayceesrk commented on 21 Dec 2021 • edited 👻	Member 😳 ···	Reviewers	ණ		
	This DD adds support for shared memory percellations through demains and			gez 🖓		
direct-style concurrency through effect handlers (without syntactic support). It intends to have backwards compatibility in terms of language features, C API, and also the performance of single-threaded code.			🍘 avsm	\Box		
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	For users		👰 sadiqj	\Box		
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• Many features broken and are being added back



This will continue after 5.0 gets released

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- Platform support
 - ✤ 32-bit will be bytecode only
 - ✤ On 64-bit,
 - x86-64 + Linux, macOS, Windows, OpenBSD, FreeBSD
 - Arm64 + Linux, macOS (Apple Silicon)
 - RISC-V (PR open)
 - ◆ JavaScript (jsoo) effect handlers are not supported yet!



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- Decoupling major slice from minor GC
- Mark stack prefetching
- Best-fit vs multicore allocator

• GC performance improvements TBD

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 - Decoupling major slice from minor GC
 - Mark stack prefetching
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- Statmemprof
 - Work in progress for reinstating asynchronous action safety



Tidying

- We tidied up accumulated deprecations
 - String.uppercase, lowercase, capitalize, uncapitalize
 - + Stream, Genlex ~> camlp-streams
 - Pervasives, ThreadUnix modules deleted
- Major version jump to make good changes
 - C function names are all prefixed uniformly
 - Additional libraries Unix, Str installed as findlib packages

OPAM Health Check

OPAM Health Check

	4.14	5.0+alpha-repo	number of revdeps
0install.2.18		\boxtimes	1
BetterErrors.0.0.1		\mathbf{X}	7
TCSLib.0.3		\mathbf{X}	1
absolute.0.1		\boxtimes	0
acgtk.1.5.3			0
advi.2.0.0		\mathbf{X}	0
aez.0.3			0
ahrocksdb.0.2.2			0
aio.0.0.3			0
alt-ergo-free.2.2.0			7
amqp-client-async.2.2.2		\mathbf{X}	0
amqp-client-lwt.2.2.2		\mathbf{X}	0
ancient.0.9.1			0
apron.v0.9.13			17



http://check.ocamllabs.io/

Sandmark Nightly Service



sandmark.tarides.com



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Sandmark Nightly Service



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OCaml 5.0 needs you!



- OCaml 4 will have longer term support than usual
- Even if you don't plan to use concurrency and parallelism, *switch to OCaml 5.0*
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 - Test, deploy, evaluate, benchmark sequential programs in 5.0
 - Report bugs & performance regressions

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 - Test, deploy, evaluate, benchmark sequential programs in 5.0
 - Report bugs & performance regressions
- **OCaml 5.0?**
 - Let us know so that we can work on it!

• What is stopping you from switching to